

# Pure\_I

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## Contents

**theory** *Pure-I* **imports** *Pure* **begin**

**typedecl** *bool*

**judgment** *Trueprop* ::  $\langle \text{bool} \Rightarrow \text{prop} \rangle$  ( $\langle \cdot \rangle$ )

**axiomatization** *Imp* (**infixr**  $\langle \longrightarrow \rangle$  3)

**where** *Imp-I* [*intro*]:  $\langle p \Longrightarrow q \Longrightarrow p \longrightarrow q \rangle$

**and** *Imp-E* [*elim*]:  $\langle p \longrightarrow q \Longrightarrow p \Longrightarrow q \rangle$

**axiomatization** *Dis* (**infixr**  $\langle \vee \rangle$  4)

**where** *Dis-E* [*elim*]:  $\langle p \vee q \Longrightarrow (p \Longrightarrow r) \Longrightarrow (q \Longrightarrow r) \Longrightarrow r \rangle$

**and** *Dis-I1* [*intro*]:  $\langle p \Longrightarrow p \vee q \rangle$

**and** *Dis-I2* [*intro*]:  $\langle q \Longrightarrow p \vee q \rangle$

**axiomatization** *Con* (**infixr**  $\langle \wedge \rangle$  5)

**where** *Con-I* [*intro*]:  $\langle p \Longrightarrow q \Longrightarrow p \wedge q \rangle$

**and** *Con-E1* [*elim*]:  $\langle p \wedge q \Longrightarrow p \rangle$

**and** *Con-E2* [*elim*]:  $\langle p \wedge q \Longrightarrow q \rangle$

**axiomatization** *Falsity* ( $\langle \perp \rangle$ )

**where** *Falsity-E* [*elim*]:  $\langle \perp \Longrightarrow p \rangle$

**definition** *Truth* ( $\langle \top \rangle$ ) **where**  $\langle \top \equiv \perp \longrightarrow \perp \rangle$

**theorem** *Truth-I* [*intro*]:  $\langle \top \rangle$

**unfolding** *Truth-def* ..

**definition** *Neg* ( $\langle \neg \rightarrow [6] \ 6 \rangle$ ) **where**  $\langle \neg p \equiv p \longrightarrow \perp \rangle$

**theorem** *Neg-I* [*intro*]:  $\langle (p \Longrightarrow \perp) \Longrightarrow \neg p \rangle$

**unfolding** *Neg-def* ..

**theorem** *Neg-E* [*elim*]:  $\langle \neg p \Longrightarrow p \Longrightarrow q \rangle$

**unfolding** *Neg-def*

**proof** –

**assume**  $\langle p \longrightarrow \perp \rangle$  **and**  $\langle p \rangle$

**then have**  $\langle \perp \rangle$  ..

**then show**  $\langle q \rangle$  ..

**qed**

**definition** *Iff* (**infixr**  $\langle \longleftrightarrow \rangle$  3) **where**  $\langle p \longleftrightarrow q \equiv (p \longrightarrow q) \wedge (q \longrightarrow p) \rangle$

**theorem** *Iff-I* [*intro*]:  $\langle (p \Longrightarrow q) \Longrightarrow (q \Longrightarrow p) \Longrightarrow p \longleftrightarrow q \rangle$

**unfolding** *Iff-def*

**proof** –

**assume**  $\langle p \Longrightarrow q \rangle$  **and**  $\langle q \Longrightarrow p \rangle$

from  $\langle p \implies q \rangle$  have  $\langle p \longrightarrow q \rangle$  ..  
 from  $\langle q \implies p \rangle$  have  $\langle q \longrightarrow p \rangle$  ..  
 from  $\langle p \longrightarrow q \rangle$  and  $\langle q \longrightarrow p \rangle$  show  $\langle (p \longrightarrow q) \wedge (q \longrightarrow p) \rangle$  ..  
 qed

theorem *Iff-E1* [*elim*]:  $\langle p \longleftrightarrow q \implies p \implies q \rangle$   
 unfolding *Iff-def*  
 proof –  
 assume  $\langle (p \longrightarrow q) \wedge (q \longrightarrow p) \rangle$   
 then have  $\langle p \longrightarrow q \rangle$  ..  
 then show  $\langle p \implies q \rangle$  ..  
 qed

theorem *Iff-E2* [*elim*]:  $\langle p \longleftrightarrow q \implies q \implies p \rangle$   
 unfolding *Iff-def*  
 proof –  
 assume  $\langle (p \longrightarrow q) \wedge (q \longrightarrow p) \rangle$   
 then have  $\langle q \longrightarrow p \rangle$  ..  
 then show  $\langle q \implies p \rangle$  ..  
 qed

proposition  $\langle p \longrightarrow \neg \neg p \rangle$   
 proof  
 assume  $\langle p \rangle$   
 show  $\langle \neg \neg p \rangle$   
 proof  
 assume  $\langle \neg p \rangle$   
 from  $\langle \neg p \rangle$  and  $\langle p \rangle$  show  $\langle \bot \rangle$  ..  
 qed  
 qed

proposition  $\langle (p \longrightarrow q) \wedge \neg q \longrightarrow \neg p \rangle$   
 proof  
 assume  $\langle (p \longrightarrow q) \wedge \neg q \rangle$   
 show  $\langle \neg p \rangle$   
 proof  
 assume  $\langle p \rangle$   
 from  $\langle (p \longrightarrow q) \wedge \neg q \rangle$  have  $\langle p \longrightarrow q \rangle$  ..  
 from  $\langle p \longrightarrow q \rangle$  and  $\langle p \rangle$  have  $\langle q \rangle$  ..  
 from  $\langle (p \longrightarrow q) \wedge \neg q \rangle$  have  $\langle \neg q \rangle$  ..  
 from  $\langle \neg q \rangle$  and  $\langle q \rangle$  show  $\langle \bot \rangle$  ..  
 qed  
 qed

proposition  $\langle (p \longleftrightarrow q) \longleftrightarrow q \longleftrightarrow p \rangle$   
 proof  
 assume  $\langle p \longleftrightarrow q \rangle$   
 show  $\langle q \longleftrightarrow p \rangle$   
 proof  
 from  $\langle p \longleftrightarrow q \rangle$  show  $\langle q \implies p \rangle$  ..  
 next  
 from  $\langle p \longleftrightarrow q \rangle$  show  $\langle p \implies q \rangle$  ..  
 qed  
 next  
 assume  $\langle q \longleftrightarrow p \rangle$   
 show  $\langle p \longleftrightarrow q \rangle$   
 proof  
 from  $\langle q \longleftrightarrow p \rangle$  show  $\langle p \implies q \rangle$  ..  
 next  
 from  $\langle q \longleftrightarrow p \rangle$  show  $\langle q \implies p \rangle$  ..

qed  
qed  
end